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TITLE OF THE INVENTION

DATA PROCESSING DEVICE

FIELD OF THE INVENTION

5 The present invention relates to a data processing device for properly scheduling conditional operation instructions in a program sequence.

DISCUSSION OF BACKGROUND

10 Conventionally, in a microprocessor for executing sequentially a plurality of instructions by pipelines, execution of branch instruction is one reason for deteriorating processing efficiency of the processor because it disturbs the pipelines. In order to prevent the processing performance from deteriorating, for
15 example, a technique of using a delay slot is proposed.

 A case that a programmer makes the following instruction sequence 1 will be described.

[Instruction sequence 1]

 (I0) CMPEQ r10,r11
20 (I1) ADD r1,r2
 (I2) BRA F0=1 H'1000
 (I3)
 (I4)

 An instruction I0 is a comparison instruction for
25 setting a flag F0 if a register 10 and a register 11 are equal as a result of a comparison therebetween. An instruction I1 is an add instruction for writing the

addition of a content of register r1 and a content of register r2 in the register r1. An instruction I2 is a conditional branch instruction for branching to an instruction in an address number 1000 of the memory.

5 Instructions I3 and I4 are arbitrary instructions and already input in the microprocessor at a time for executing the instruction I2. When a branch is taken as a result of execution of the instruction I2, the instructions I3 and I4 under the pipeline processing are
10 invalidated.

Accordingly, in consideration of the instruction I1, which is an instruction executed regardless of the branch by the instruction I2 and not an instruction of performing an operation of determining a branch
15 condition of the instruction I2, a scheduling of the instructions as shown in the following instruction sequence 2.

[Instruction sequence 2]

(I0) CMPEQ r10,r11
20 (I2) BRA F0=1 H'1000
(I1) ADD r1,r2
(I5) NOP

Even though a branch to the address number 1000 is determined as a result of the execution of instruction
25 I2, the instruction I1, which is inputted into the pipeline and processed, can further be executed without invalidating the same, wherein in a case of architecture

introducing two instructions into the pipeline at the time of executing the instruction I2, a delay slot is occupied by these two instructions. In the case of instruction sequence 2, a delay slot is constituted by the instruction I1 and the instruction I5. The instruction I5 is a so-called no operation (NOP) instruction. After the branch is taken by the instruction I2, an instruction after this branch will be fetched after the instruction I5.

Such a scheduling from the instruction sequence 1 to the instruction sequence 2 is performed by a program itself or a compiler.

Concerning such a scheduling of instructions, various techniques were proposed in "Computer Architecture: A Quantitative Approach, Morgan Kaufmann Co., year 1990".

Thus, a technology of scheduling instructions in a program was important in order to draw out a processing capability of a microprocessor as much as possible. However, there were various restrictions on the scheduling depending on a type of instruction. In the above case, the conditional branch instruction by the instruction I2 could not be posed before the instruction I0 determining the execution condition of the conditional branch instruction in their instruction sequence. This was because the execution of instruction I0 of determining the branch condition before

instruction I2 referred to a content of the flag F0 which was the branch condition by the instruction I2. Thus, not only a conditional branch instruction but also a conditional arithmetic operation instruction, were
5 reasons for deteriorating a degree of freedom in the scheduling of instruction.

In a microprocessor of a type of very long instruction word (VLIW), a plurality of instructions, which could be executed in parallel, were expressed by a
10 single instruction set. In this type, it was necessary to use a high-level scheduling technology considering that which instructions could be executed in parallel, whereby many circumstances that instructions which were meaningless in respect of the program, namely, so-called
15 no operation (NOP) instructions, were inserted in an instruction sequence occurred because of existence of conditional operation instructions. It was also a reason for deteriorating a processing performance of microprocessor to process an no operation (NOP)
20 instruction.

SUMMARY OF THE INVENTION

It is an object of the present invention to solve the above-mentioned problems inherent in the prior art and to provide a circumstance under which a scheduling
25 of instructions having a high degree of freedom is obtainable in a data processing for processing conditional operation instructions for a programmer.

According to a first aspect of the present invention, there is provided a data processing device comprising an instruction decoder for outputting a control signal corresponding to each operation
5 instruction by successively decoding a plurality of coded operation instructions described in a program sequence and an instruction execution unit for executing operations which are respectively designated by the plurality of operation instructions in accordance with
10 the control signals outputted from the instruction decoder, wherein a first operation instruction is decoded in a first period and the operation designated thereby is executed in a second period following the first period. Meanwhile, a second operation instruction,
15 of which operation is executed under a predetermined condition, is decoded in a third period, it is judged whether the predetermined condition is satisfied in a fourth period, which is started after the same time as the second period or a longer time than the second
20 period from the ending of the third period, and the instruction execution unit executes an operation designated by the second operation unit in response to a result of the judgment.

The data processing device further comprises a
25 register for designating amount of delay that can variably set a hold value. The instruction execution unit starts a judgement of whether or not the second

operation instruction satisfies the predetermined condition in response to a value held as an amount of delay in the register for designating amount of delay.

The coded second operation instruction has a field
5 for designating operation and a field for designating amount of delay that designates an interval between the ending of the third period and the starting of the fourth period, wherein the amount of delay is set in the register for designating amount of delay in accordance
10 with a content described in the field for designating amount of delay.

The data processing device further comprises a program counter for sequentially counting an address corresponding to each of a plurality of operation
15 instructions and holding the address. The register for designating amount of delay is to hold an address value as an amount of delay. The instruction execution unit starts to judge whether or not the second operation instruction satisfies the predetermined condition in
20 response to an event that the address value held in the register for designating amount of delay is in agreement with the value in the program counter.

The instruction execution unit judges whether or not the predetermined condition is satisfied in a fifth
25 period included in the fourth period. In this, the instruction execution unit executes an operation designated by the second operation instruction when the

predetermined condition is satisfied in a sixth period included in the fourth period and starting after passing the same time as the second period or a longer time than the second period from the ending of the fifth period.

5 The data processing device comprises a first register for designating amount of delay and a second register for designating amount of delay, in both of which respective hold values can be set variably. The instruction execution unit starts a judgement of whether
10 or not the second operation instruction satisfies the predetermined condition in accordance with the hold value as a first amount of delay in the first register for designating amount of delay and executes by starting the operation designated by the second operation
15 instruction when the second operation instruction satisfies the predetermined condition in accordance with the hold value held as a second amount of delay in the second register for designating amount of delay.

 The coded second operation instruction comprises a
20 field for designating operation, a first field for designating amount of delay which designates a time between the ending of the third period and the starting of the fourth period, and a second field for designating amount of delay which designates a time from the ending
25 of the fifth period and the starting of the sixth period. The first register for designating amount of delay is set with a first amount of delay in accordance

with a content described in the first field for
designating amount of delay; and the second register for
designating amount of delay is set with a second amount
of delay in accordance with a content described in the
5 second field for designating amount of delay.

The first register for designating amount of delay
and the second register for designating amount of delay
hold values of address respectively as the first amount
of delay and the second amount of delay. The instruction
10 execution unit starts a judgement of whether or not the
predetermined condition is satisfied in response to an
event that the value of address held in the first
register for designating amount of delay is in agreement
with a value of program counter. Further, the
15 instruction execution unit starts to execute the
operation designated by the second operation instruction
when the predetermined condition is satisfied in
response to an event that the value of address held in
the second register for designating amount of delay is
20 in agreement with the value of program counter.

With respect to a third operation instruction among
a plurality of operation instructions, the instruction
decoder decodes in a seventh period which starts after
the third period, and an instruction execution unit
25 designated by the third operation instruction in an
eighth period, which starts after the seventh period,
executes an operation; and a result of the operation is

written in a predetermined memory location. At this time, the second operation instruction designates an operation instruction which would be executed in a case that the result of operation by the third operation

5 instruction has a predetermined value; and the instruction execution unit determines whether or not the operation is executed in reference of the predetermined memory location so that the starting of the fourth period becomes at least after the eighth period.

10 The predetermined memory location is a flag or a register. The third operation instruction is a comparison instruction for comparing values of the two registers and writing a result of the comparison in the predetermined memory location. In this, the second
15 operation instruction is a branch instruction, a jump instruction or an add instruction.

Each of the plurality of operation instructions has a field for designating operation that designates a content of the operation, a field for designating
20 condition that designates an execution condition of the operation and a field for designating amount of delay that designates an amount by which timing for judging the execution condition is delayed.

A description of representing that the field for
25 designating condition of the first operation instruction is unconditional is described provided that the first operation instruction is an instruction executable

unconditionally. The instruction decoder outputs a first control signal in accordance with the field for designating operation in the first operation instruction and controls the instruction execution unit so as to
5 execute the operation designated by the first operation instruction in the second period.

Further, when the first operation instruction is an unconditional operation instruction, a description representing a condition other than the predetermined
10 condition of the second operation instruction and a description of judging said other condition in the first period are described respectively in the field for designating condition and the field for designating amount of delay.

15 The instruction decoder outputs the first control signal in accordance with the field for designating operation in the first operation instruction and controls the instruction execution unit to execute the operation designated by the first operation instruction
20 in the second period based on the field for designating condition and the field for designating amount of delay.

The instruction decoder judges whether or not the condition is satisfied in accordance with the field for designating condition and the field for designating
25 amount of delay in the first operation instruction to thereby output the first control signal in accordance with the field for designating operation in the first

operation instruction in response to a result of the judgement.

In a field for designating condition and a field for designating amount of delay in the second operation
5 instruction, a description representing predetermined conditions and a description representing an interval between the ending of the third period and the starting of the fourth period are described. The instruction decoder outputs the first control signal in accordance
10 with the field for designating operation in the second operation instruction; controls the instruction execution unit to judge whether or not the predetermined condition is satisfied in the fourth period in accordance with the field for designating amount of
15 delay in the second operation instruction; and controls the instruction execution unit to determine whether or not the predetermined condition is satisfied in accordance with the field for designating condition in the second operation instruction.

20 According to a second aspect of the present invention, there is provided a data processing device comprising an instruction decoder for decoding a conditional operation instruction and outputting a first control signal and an instruction execution unit for
25 executing an operation in accordance with the first control signal. The instruction execution unit includes a first register for holding the first control signal, a

second register for holding a first description
representing a condition of executing an operation
designated by the conditional operation instruction and
a third register for holding a second description
5 representing a time for starting a judgement of the
condition. The instruction execution unit starts to
judge whether or not the condition is satisfied based on
the first description held in the second register in
response to an event that the time for starting the
10 judgement of the condition is detected based on the
second description held in the third register, and
starts to execute the operation designated by the
operation instruction after reading out the first
control signal held in the first register in response to
15 a result of the judgement.

The second description held in the third register
can be set variably.

The data processing device has a program counter for
successively counting addresses respectively
20 corresponding to a plurality of operation instructions
and holding the addresses. The third register holds a
value of address as the second description; the
instruction execution unit detects an event that the
value of address held in the third register is in
25 agreement with the address in the program counter; and
starts a judgement of whether or not the condition is
satisfied in response to the detection.

The conditional operation instruction includes a field for designating operation which designates a content of operation, a field for designating condition which designates an execution condition of the operation, and a field for designating amount of delay which designates a time for judging the execution condition. The instruction decoder produces the first control signal based on a content described in the field for designating operation and outputs a content described in the field for designating condition as the first description and outputs a content described in the field for designating amount of delay. The first description outputted from the instruction decoder is held in the second register. The instruction execution unit writes the second description in the third register in accordance with the field for designating amount of delay outputted from the instruction decoder.

The instruction execution unit further includes a fourth register for holding a third description representing a time for starting an operation designated by the operation instruction. The instruction execution unit detects the time for starting the operation designated by the operation instruction in accordance with the third description; judges whether or not the condition is satisfied in response to a result of the detection; and starts to execute the operation designated by the operation instruction in response to

the result of judgement.

The data processing device has a program counter for successively counting addresses respectively corresponding to a plurality of instructions and holding
5 the same. A value of address is held in the third register as the second description. A value of address different from the second description is held in a fourth register. The instruction execution unit detects an event that the value of address held in the third
10 register is in agreement with the address in the program counter; and starts a judgement of whether or not the condition is satisfied in response to the detection. Further, the instruction execution unit detects an event that the value of address held in the fourth register is
15 in agreement with the address in the program counter; and starts to execute an operation designated by the operation instruction in response to the detection.

The coded conditional operation instruction includes a field for designating operation which designates a
20 content of operation, a field for designating condition which designates an execution condition of the operation, a first field for designating amount of delay which designates a time for judging the execution condition and a second field for designating amount of
25 delay which designates a time for starting execution of the operation. The instruction decoder produces the first control signal based on a content described in the

field for designating operation; outputs a content described in the field for designating condition as the first description; and outputs contents described in the first field for designating amount of delay and the
5 second field for designating amount of delay. The instruction execution unit writes the second description in the third register in accordance with the content described in the first field for designating amount of delay outputted from the instruction decoder and writes
10 the third description in the fourth register in accordance with the content described in the second field for designating amount of delay outputted from the instruction decoder.

BRIEF DESCRIPTION OF THE DRAWINGS

15 A more complete appreciation of the invention and many of the attendant advantages thereof will be readily obtained as the same becomes better understood by reference to the following detailed description when considered in connection with the accompanying drawings,
20 wherein:

Figure 1 is a block chart for showing a constitution of a microprocessor according to Embodiment 1 of the present invention;

Figure 2a is a chart for schematically showing an
25 instruction format of the microprocessor;

Figure 2b is a chart for schematically showing an instruction format of the microprocessor;

Figure 3a is a chart for schematically showing a pipeline operation at a time of executing parallel sub-instructions in the microprocessor;

Figure 3b is a chart for schematically showing a
5 pipeline operation at a time of executing parallel sub-instructions in the microprocessor;

Figure 3c is a chart for schematically showing a pipeline operation at a time of executing parallel sub-instructions in the microprocessor;

10 Figure 4 is a chart for showing detailed contents of operation fields;

Figure 5a is a chart for showing a constitution of register of the microprocessor;

Figure 5b is a chart for showing a constitution of
15 register of the microprocessor;

Figure 5c is a chart for showing a constitution of register of the microprocessor;

Figure 6 is a chart for showing a detailed contents of PSW;

20 Figure 7 is a chart for showing a basic format of delayed branch sub-instruction;

Figure 8 shows an example of program which is processed in the microprocessor shown in Figure 1 for explanation;

25 Figure 9 is a chart for schematically showing a pipeline operation of the microprocessor at a time of processing the program shown in Figure 8;

Figure 10 shows an example of another program processed in the microprocessor shown in Figure 1;

Figure 11 is a chart for schematically showing a pipeline operation in the microprocessor at a time of
5 processing the program shown in Figure 9; and

Figure 12 is a block chart for schematically showing a constitution of the microprocessor according to Embodiment 2 of the present invention.

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENTS

10 A detailed explanation will be given of preferred embodiment of the present invention in reference to Figures 1 through 12 as follows, wherein the same numerical references are used for the same or the similar portions and description of these portions is
15 omitted.

EMBODIMENT 1

Figure 1 is a block chart for showing a constitution of a microprocessor according to Embodiment 1 of the present invention. This microprocessor is a 32-bit
20 microprocessor having an internal data bus of 32 bits. In this Figure, reference numeral 2 designates an instruction decode unit (instruction decoder) for performing a process of decoding an instruction code inputted from an instruction RAM 6 through an ID bus
25 having a width of 64 bits; reference numeral 3 designates a memory unit (instruction execution unit) for performing an address calculation; reference numeral

4 designates an integer operation unit (instruction execution unit) for performing an arithmetic logic operation and/or a shift operation; reference numeral 5 designates a general purpose register of 32 bits \times 64 words; and reference numeral 7 designates a data RAM for storing data.

In the instruction decode unit 2, reference numerals 8 and 9 respectively designate decoders which decode an instruction code; and reference numeral 10 designates a processor status word (hereinbelow, referred to as PSW) which represents a status of the processor. The instruction decode unit 2 makes a control signal 11 based on a result of decoding in the decoder 8 and a content of PSW 10, and transfers the control signal 11 to the memory unit 3. On the other hand, the instruction decode unit 2 makes a control signal 12 based on a result of decoding in the decoder 9 and a content of PSW 10, and transfers the control signal 12 to the integer operation unit 4.

In the memory unit 3, reference numeral 13 designates a PC controlling part which calculates a program counter value (PC value) with respect to an instruction to be executed subsequently by adding 8 to the PC value when a sub-instruction excluding jump and branch is executed; adds a branch displacement to the PC value when a sub-instruction including jump and/or branch is executed; and calculates the PC value with

respect to an instruction existing in a destination of
jump by calculating in consideration of an addressing
mode designated by an operation. Also, the PC
controlling part 13 transfers the calculated PC value to
5 the instruction RAM 6 through an IA bus having a width
of 32 bits so that an instruction code is outputted from
the instruction RAM 6. Reference numeral 14 designates a
memory controlling part for controlling accesses to data
which are to be an operand. The memory controlling part
10 14 transfers address data to the data RAM 7 through a DA
bus having a width of 32 bits and accesses to data
necessary for instruction execution through a DD bus
having a width of 64 bits. Reference numeral 15
designates an ALU which performs an arithmetic logic
15 operation using data having maximum three words
transferred from the general purpose register 5 through
a S1 bus, S2 bus, and S3 bus each having a width of 32
bits and transfers a result of operation to the general
purpose register 5 through the D1 bus having a width of
20 32 bits; and reference numeral 16 designates a shifter
which performs a shift operation using data transferred
from the general purpose register 5 through the S1 bus,
S2 bus and S3 bus and transfers a result of operation to
the general purpose register 5 through the D1 bus.

25 It is possible to transfer data having a length of
32 bits as much as 4 words simultaneously to the memory
unit 3 through the S1 bus, the S2 bus, the S3 bus and an

S4 bus. Accordingly, it is possible to realize a two-word store sub-instruction which, for example, stores a memory area addressed by the sum of a content of first register and a content of second register with a content of third register and simultaneously stores a memory area addressed by a value obtained by adding a predetermined value to an address storing the content of third register with a content of fourth register. Also, it is possible to transfer a result of two words obtained by an operation in the memory unit 3 or two-word data transferred from the data RAM 7 to the general purpose register 5 through the D1 bus and the D2 bus.

The PC controlling part 13 includes registers 30, 31, 32, and 33 and a memory circuit 34 for holding one bit. The memory controlling part 14 includes registers 40, 41, and 42. The ALU 15 includes registers 50, 51 and 52. The shifter 16 includes registers 60, 61, and 62. These registers will be described in the below.

In the integer operation unit 4, reference numeral 17 designates a multiplier which performs a multiplication using data having maximum three words transferred from the general purpose register 5 through S4 bus, S5 bus, and S6 bus each having a width of 32 bits and transfers a result of the operation to the general purpose register 5 through D2 bus and D3 bus each having a width of 32 bits; and reference numeral 18 designates an accumulator which holds by cumulatively

adding or cumulatively subtracting the result of multiplication. As for the accumulator, there are provided a pair of 64-bit accumulators. Reference numeral 19 designates an ALU which performs an

5 arithmetic logic operation using data having maximum three words transferred from the general purpose register 5 through S4 bus, S5 bus and S6 bus and transfers a result of the operation to the general purpose register 5 through D2 bus and D3 bus; and

10 reference numeral 20 designates a shifter which performs a shift operation using data transferred from the general purpose register 5 through S4 bus, S5 bus and S6 bus and transfers a result of the operation to the general purpose register 5 through D2 bus and D3 bus.

15 The multiplier 17 includes registers 70, 71, and 72. The ALU 19 includes registers 80, 81, and 82. The shifter 20 includes registers 90, 91, and 92. These registers will also be described in the below.

This microprocessor can read out register values of

20 maximum 6 kinds from the general purpose register 5, wherein the read-out data are outputted respectively to S1 bus, S2 bus, S3 bus, S4 bus, S5 bus, and S6 bus. Also, register values of maximum three kinds can be written in the general purpose register 5 simultaneously

25 through D1 bus, D2 bus, and D3 bus.

Figures 2a and 2b are views for explaining instruction formats of the microprocessor 1. In the

instruction formats, there are included a format 101 of dual operation instruction which designates two operations by one instruction word as shown in Figure 2a and a format 102 of single operation instruction which
5 designates one operation by one instruction word as shown in Figure 2b.

In the format 101 of dual operation instruction, there is included a format field composed of a field 103 and a field 104, two operation fields 106 and 107,
10 execution condition fields 401 and 402 respectively attached to the operation fields 106 and 107, and fields for designating amount of delay for judging condition (hereinbelow, referred to as CD field) 404 and 405
respectively attached to the execution condition fields
15 401 and 402.

The format 102 of single instruction operation 102 includes a format field composed of fields 103 and 104, an operation field composed of fields 108, 109, and 110, an execution condition field 403 attached to the
20 operation field and a CD field 406 attached to the execution condition field 403.

The format fields have the following meanings.

Code: format	Order of execution	
	operation_0	operation_1
FM= 00:2 sub-instructions	first	first
01:2 sub-instructions	first	second

10:2	sub-instructions	second	first
11:1	one sub-instruction	first	--

In this, reference FM has a value of 2 bits composed
5 of the field 103 and the field 104.

A plurality of pipeline stages in the microprocessor
1 are formed of an instruction fetch stage IF, an
instruction decode stage D/A, an instruction execution
stage E/M and a write back stage W, wherein processings
10 in each stage are finished within one clock cycle. Figure
3a through 3c are schematic views for explaining the
pipeline stages for processing the dual operation
instruction 101 in the microprocessor 1. In the
instruction fetch unit IF, the dual operation
15 instruction 101 is fetched from a memory RAM 6 to an
instruction decode unit 2. In the instruction decode
stage D/A, operation_0 described in the operation field
106 is decoded by the decoder 8 and operation_1
described in the operation field 107 is decoded by the
20 decoder 9. Further, addresses of each operand of
operation_0 and operation_1 or, in a case that operation
0 and operation_1 are branch sub-instructions, addresses
of destination of branch are calculated in the
instruction decode stage D/A. In the instruction
25 execution stage E/M, an operation designated by
operation_0 in accordance with a control signal 11 is
executed in the memory unit 3, and an operation

designated by operation_1 in accordance with a control signal 12 is executed in the integer operation unit 4. When operation_1 is a sub-instruction accompanying memory access such as a load sub-instruction and a store sub-instruction, the memory unit 3 accesses to the memory in the instruction execution stage E/M. In the write back stage W, a result of operation obtained in the memory unit 3 and a result of operation obtained in the integer operation unit 4 are written in registers designated by operation_0 and operation_1 respectively. In a sub-instruction which does not accompany a sub-instruction of writing a result of operation in a register of the processor 1, namely, a branch sub-instruction, a jump sub-instruction, a store sub-instruction storing a memory with data, a comparison sub-instruction reflecting a result of operation in a flag and so on, a write back stage W is not included. Depending on a microprocessor, a write back stage W is processed in the same clock cycle as that of an instruction execution stage E/M.

In the case of FM=00, the stages IF, D/A, E/M, and W respectively of operation_0 and operation_1 are performed in parallel, and operation_0 and operation_1 are processed in 4 clocks, as shown in Figure 3a.

In the case of FM=01, the stages IF, D/A, E/M and W of operation_0 are continuously performed in 4 clocks, as shown in Figure 3b. The stages IF and D/A

respectively of operation_0 and operation_1 are performed in parallel. Meanwhile, the stages E/M and W of operation_1 are performed with a delay of 1 clock from those of operation_0. The stage E/M of operation_1 is performed in parallel with the stage W of operation_0. Accordingly, operation_1 is processed in 5 clocks. In the case of FM=01, the stages IF, D/A, E/M and W are continuously performed in 4 clocks, as shown in Figure 3c. The stages IF and D/A respectively of operation_0 and operation_1 are performed in parallel. Meanwhile, the stages E/M and W of operation_0 are performed with a delay of 1 clock from those of operation_1. The stage E/M of operation_0 is performed in parallel with the stage W of operation_1. Accordingly, operation_1 is processed in 5 clocks.

Also, a single operation instruction 102 having a format shown in Figure 2b is also processed in 1 clock cycle in each of the instruction fetch stage IF, the instruction decode stage D/A, the instruction execution stage E/M and the write back stage W. In the stage IF, the single operation instruction 102 is fetched from the instruction RAM 6 to the instruction decode unit 2. In the stage D/A, the single operation instruction 102 is inputted respectively in the decoders 8 and 9. In response to a type of operation designated by the single operation instruction 102, one of the decoders 8 and 9 decodes the single operation instruction 102. When the

decoder 8 decodes, it outputs a control signal 11, and when the decoder 9 decodes, it outputs a control signal 12. In the stage E/M, the memory unit 3 (or the integer operation unit 4) executes the operation designated by the single operation instruction 102 in accordance with the control signal 11 (or the control signal 12). In the stage W, a result of operation obtained in the stage E/M is written in a register designated by the single operation instruction 102.

10 In the next, codes of execution condition will be described. Every execution condition fields 401, 402, and 403 has the following meaning.

Code: Execution condition

15 CC=000: always
001: F0=true and F1=don't care
010: F0=false and F1=don't care
011: F0=don't care and F1=true
100: F0=don't care and F1=false
20 101: F0=true and F1=true
110: F0=true and F1=false
111: reserved

Each of the execution condition fields 401, 402, and 403 is provided to designate an execute condition of operation sub-instruction attached thereto for determining whether operations of operation_0 and operation_1 in the operation fields 106 and 107 and

operations in the operation field composed of the fields 108, 109, and 110 are valid or invalid according to execution control flags F0 and F1. The execution control flags F0 and F1 exist in a processor status word (PSW) 5 10 as described in the below. When the operation is valid, a result of the operation is reflected in a register, a memory and a flag and a result of an operation obtained by the operation is left. On the other hand, when the operation is invalid, the operation 10 designated by decoding the operation sub-instruction is not executed or a result of the operation is not reflected on a register, a memory, nor a flag even though the operation is executed, namely the no result of operation is left as if an invalid operation (NOP) is 15 executed.

When a value of execution condition field CC is 000, an operation is always valid despite values of execution control flags F0 and F1. When CC=001, an operation is valid only in a case that the execution control flag F0 20 is true. It doesn't care the state of control flag F1. When CC=010, the operation is effective only in a case that the execution control flag F0 is false. It doesn't care the state of execution control flag F0. When CC=011, the operation is effective only in a case that 25 the execution control flag F1 is true. It doesn't care the state of execution control flag F0. When CC=100, the operation is valid only in a case that the execution

control flag F1 is false. It doesn't care the state of execution control flag F0. When CC=101, the operation is valid only in a case that the execution control flag F0 is true and simultaneously the execution control flag F1 is true. When CC=110, the operation is valid in a case that the execution control flag F0 is true and the execution control flag F1 is false. When CC=111, the process is undefined and a user can not use an instruction defined by CC=111.

10 Figure 4 is a diagram for showing detailed contents of operation fields. Formats 111 through 117 are used for a short-type operation field 106 and a short-type operation field 107 both expressed by 28 bits. Format 118 is used for a long-type operation field composed of
15 fields 108, 109 and 110.

 The format 111 (Short_M) is composed of a field 120 for designating a content of operation, two fields 121 and 122 for designating register numbers, a field 123 for designating a register number of an immediate value
20 having a length of 6 bits and a field 124 for designating whether the field 123 designates a register number or an immediate value. As shown in Figures 3a through 3c, when a value X in the field 124 is "00", "01" or "11", it means that the field 123 designates a
25 register number. When the value X of field 124 is "10", it means that the field 123 designates an immediate value.

This format 111 is used for a memory access operation when a register indirect addressing is conducted.

The format 112 (Short_A) is composed of a field 120 for designating a content of operation, two fields 121 and 122 for designating register numbers, a field 123 for designating a register number or an immediate value having a length of 6 bits and a field 125 for designating whether the field 123 designates the register number or the immediate value. As shown in Figures 3a through 3c, when a value X' of field 125 is "0", it means that the field 123 designates the register number. When the value X' of field 125 is "1", it means that the field 123 designates the immediate value.

This format 112 is used for an arithmetic operation, a logical operation, a shift operation and a bit operation.

The format 113 (Short_B1) is composed of a field 120 for designating a content of operation and a field 126 for designating a register number. This format 113 is used for a jump instruction and a branch instruction based on a designation of register.

The format 114 (Short_B2) is composed of a field 120 for designating a content of operation and a field 127 having a displacement of 18-bit length. This format 114 is used for a jump instruction or a branch instruction.

Format 115 (Short_B3) is composed of a field 120 for

designating a content of operation, a field 121 for
designating a register number, a field 128 for
designating a register number or an immediate value
having a length of 12 bits, a field 129 for designating
5 whether the field 128 designates the register number or
the immediate value and a field 130 for designating
whether or not a conditional jump or a conditional
branch is executed based on the field 121 upon judgement
of 0.

10 This format 115 is used for a conditional jump
instruction or a conditional branch instruction.

Format 116 (Short_D1) is composed of a field 120 for
designating a content of operation, a field 121 for
designating a register number, a field 128 for
15 designating a register number or an immediate value
having a length of 12 bits and a field 129 for
designating whether or not the field 128 designates the
register number or the immediate value.

This format 116 is used for a conditional jump
20 instruction, a conditional branch instruction and a
repeat instruction.

The format 117 (Short_D2) is composed of a field 120
for designating a content of operation, a field 128 for
designating a register number or an immediate value
25 having a length of 12 bits, a field 129 for designating
whether the field 128 designates the register number or
the immediate value and a field 131 relating to a

delayed instruction.

This format 117 is used for a delayed jump instruction, a delayed branch instruction, and a repeat instruction.

5 The format 118 (Long) is composed of a field 120 for designating a content of operation, two fields 121 and 122 for designating register numbers and a field 132 for designating an immediate value having a length of 32 bits.

10 This format 118 is used for a complicated arithmetic operation, an arithmetic operation using a large immediate value, a memory access operation in register indirect addressing with a large displacement, a branch operation accompanying a large displacement, a jump
15 instruction to an absolute address, and so on.

Figures 5a through 5c are diagrams for explaining a register configuration of a microprocessor.

This microprocessor has general purpose registers 5 having a length of 32 bits as many as a number of 64 as
20 shown in Figure 5a, control registers 150 as many as a number of 12 as shown in Figure 5b and accumulators 18 as many as a number of 2 as shown in Figure 5c. A content 140 in a general purpose register R0 is always 0 and a writing thereto is ignored. General purpose
25 register R62 is a link register in which a return address from a subroutine is set. The general purpose register R63 is a stack pointer, which works as a user

stack pointer (SPU) or an interrupt stack pointer (SPI) in response to a value of SM field in PSW10. In the control registers 150, a program counter 151, PSW10 and various registers for dedicated use are included.

5 In an operation by the format 112 shown in Figure 4, each of upper 16 bits and lower 16 bits of the 64 general purpose registers 5 can be separately accessed. Also, each of upper 32 bits and lower 32 bits of the two accumulators 18 can be separately accessed.

10 Figure 6 is a diagram for showing a detailed content of PSW10.

Upper 16 bits of PSW10 include a SM field 171 for switching a stack pointer, an EA field for indicating detection of a self debug trap (SDBT), a DB field 173
15 for designating admission of SDBT, an IE field 174 for designating interrupt permission, an RP field 175 for designating permission of repeat operation and an MD field 176 for designating permission of a modulo addressing. Lower 16 bits are a flag field 180. The
20 flag field 180 includes eight flags, wherein an F0 flag 181 and an F1 flag 182 designate validness/invalidness of operation. A value of each flag changes depending on a result of comparison operation or a result of arithmetic operation. Further, it changes by
25 initializing by an operation of initializing flag or by writing an arbitrary value in the flag field 180 using an operation of writing flag value. A content of the

flag field 180 can be read out by an operation of reading flag value.

Each field has the following meaning.

- | | |
|----|---|
| 5 | SM=0: stack mode 0 → use SPI |
| | SM=1: stack mode 1 → use SPU |
| | EA=0: SDBT is not detected |
| | EA=1: SDBT is detected |
| | DB=0: SDBT is not enabled |
| 10 | DB=1: SDBT is enabled |
| | IE=0: interrupt is not enabled |
| | IE=1: interrupt is enabled |
| | RP=0: repeat block is invalid |
| | RP=1: repeat block is valid |
| 15 | MD=0: modulo addressing is invalid |
| | MD=1: modulo addressing is valid |
| | F0: general purpose flag (execution control flag) |
| | F1: general purpose flag (execution control flag) |
| | F2: general purpose flag |
| 20 | F3: general purpose flag |
| | F4(S): saturated operation flag |
| | F5(V): overflow flag |
| | F6(VA): cumulative overflow flag |
| | F7(C): carry/borrow flag |
| 25 | |

Hereinbelow, a sub-instruction list of this microprocessor is shown.

A. Sub-instructions concerning function of
microprocessor

A-1. Load/store sub-instruction

5	LDB:	Load one byte to a register with sign extension
	LDBU:	Load one byte to a register with zero extension
	LDH:	Load one half-word to a register with sign extension
10	LDHH:	Load one half-word to a register high
	LDHU:	Load one half-word to a register with zero extension
	LDW:	Load one word to a register
15	LD2W:	Load two words to registers
	LD4BH:	Load four bytes to four half-words in two registers with sign extension
	LD4BHU:	Load four bytes to four half-words in two registers with zero extension
20	LD2H:	Load two half-words to two words in two registers with sign extension
	STB:	Store one byte from a register
	STH:	Store one half-word from a register
	STHH:	Store one half-word from a register high
25	STW:	Store one word from a register
	ST2W:	Store two words from registers
	ST4TB:	Store four bytes from four half-words from

two registers

ST2H: Store two half-words from two registers

MODDEC: Decrement a register value by a 5-bits
immediate value

5 MODINC: Increment a register value by a 5-bits
immediate value

A-2. Transfer sub-instruction

MVFSYS: Move data from a control register to a
general purpose register

10 MVTSYS: Move data from a general purpose register to
a control register

MVFACC: Move data from an accumulator to two general
purpose registers

MVTACC: Move data from two general purpose registers
15 to an accumulator

A-3. Comparison instruction

CMPcc: Compare
cc=EQ (equal), NE(not equal), GT(greater
than),

20 GE(greater than or equal to), LT(less than),
LE(less than or equal to), PS(both are
plus),

NG(both are minus)

CMPUcc: Compare unsigned

25 cc= GT, GE, LT, LE

A-4. Maximum value/minimum value sub-instruction

reserved

A-5. Arithmetic operation sub-instruction

	ABS:	Absolute
	ADD:	Add
	ADDC:	Add with carry
5	ADDHppp:	Add half-word ppp=LLL (lower of register, lower of register, lower of register), LLH (lower of register, lower of register, upper of register), LHL, LHH, HLL, HLH, HHL, HHH
10	ADDS:	Add register Rb with the sign of the third operand
	ADDS2H:	Add sign to two half-words
	ADD2H:	Add two pairs of half-words
15	AVG:	Average with rounding towards positive infinity
	ADG2H:	Average two pairs of half-words respectively rounding towards positive infinity
	JOINpp:	Join two half-words pp=LL, LH, HL, HH
20	SUB:	Subtract
	SUBB:	Subtract with borrow
	SUBHppp:	Subtract half-word ppp=LLL, LLH, LHL, LHH, HLL, HLH, HHL, HHH
25	SUB2H:	Subtract two pairs of half-words

A-6. Logical operation sub-instruction

AND: logical AND
OR: logical OR
NOT: logical NOT
XOR: logical exclusive OR
5 ANDFG: logical AND flags
ORFG: logical OR flags
NOTFG: logical NOT a flag
XORFG: logical exclusive OR flags

A-7. Shift operation sub-instruction

10 SRA: Shift right arithmetic
SRA2H: Shift right arithmetic two half-words
SRC: Shift right concatenated registers
SRL: Shift right logical
SRL2H: Shift right logical two half-words
15 ROT: Rotate right
ROT2H: Rotate right two half-words

A-8. Bit operation sub-instruction

BCLR: Clear a bit
BNOT: Invert a bit
20 BSET: Set a bit
BTST: Test a bit

A-9. Branch sub-instruction

BRA: Branch
BRATZR: Branch if zero
25 BRATNZ: Branch if not zero
BSR: Branch to subroutine
BSRTZR: Branch to subroutine if zero

BSRTNZ: Branch to subroutine if not zero

JMP: Unconditional jump

JMPTZR: Jump if zero

JMPTNZ: Jump if not zero

5 JSR: Jump to subroutine

JSRTZR: Jump to subroutine if zero

JSRTNR: Jump to subroutine if not zero

NOP: No operation

10 Delayed branch, jump sub-instruction

DBRA: Delayed branch

DBRAI: Delayed branch immediate (immediate value)

DBSR: Delayed branch to subroutine

DBSRI: Delayed branch immediate to subroutine
(immediate value)

15

DJMP: Delayed jump

DJMPI: Delayed jump immediate (immediate value)

DJSR: Delayed jump to subroutine

DJSRI: Delayed jump immediate to subroutine

20 (immediate value)

A-10. Sub-instruction concerning OS

TRAP: Trap

REIT: Return from exception, interrupts, and traps

B. Sub-instruction concerning DSP function

25 B-1. Arithmetic operation sub-instruction

MUL: Multiply

MULX: Multiply with extended precision (double

precision)

MULXS: Multiply and shift to the right by one with
extended precision (double precision)

MULX2H: Multiply two pairs of half-words with
extended precision (double precision)

MULHXpp: Multiply two half-words with extended
precision (double precision)
pp=LL, LH, HL, HH

MUL2H: Multiply two pairs of half-words

MACa: Multiply and add (operation)
a (designating accumulator) = 0, 1

MACSa: Multiply, shift to the right by one and add
(operation)
a=0, 1

MSUBa: Multiply and subtract (operation)
a=0, 1

MSUBSa: Multiply, shift to the right by one and
subtract (operation)
a=0, 1

B-2. Repeat sub-instruction

REPEAT: Repeat a block of instructions

REPEATI: Repeat a block of instructions immediate
(designating immediate value)

In the next, the CD field in the instruction formats
shown in Figures 2a and 2b will be described in detail.
The CD field 404 is to designate an amount of delay by

which timing of judging an execution condition
designated by the execution condition field 401
corresponding thereto in a pipeline process of
operation_0 designated by the operation field 106

- 5 corresponding thereto, wherein the amount of delay can
be variably set when a user sets a value of CD field 404
appropriately.

- Specifically, in the CD field 404, an offset value
OVA from a memory address X of the instruction with the
10 format 101 is described as an immediate value. In this
case, timing of judging the execution condition
described in the execution condition field 401 in the
process of operation_0 is a clock cycle when a PC value
of the microprocessor 1 holds the address No. (X+OVA),
15 wherein the offset value OVA can be 0. In such a case,
the timing of judging the execution condition is a clock
cycle when the PC value holds the address No. X.

- Further, it is possible to describe a description
designating a register number of a register in the
20 processor 1 which holds an address value in the CD field
404. In this case, the timing of judging the execution
condition described in the execution condition field 401
is a clock cycle where the PC value of the
microprocessor 1 is an address held in the designated
25 register. A bit for distinguishing whether an immediate
value is described in the CD field 404 or a register
number is described therein exists in the identical CD

field 404.

However, when the execution condition field 401 designates the unconditional execution of CC=000, the value of CD field 404 is ignored by the instruction
5 decode unit 2 and the units 3 and 4 at the process of executing operation_0 corresponding thereto.

The CD field 405 works the same as the CD field 404 does with respect to the operation field 107 and the condition execution field 402. Further, the CD field 406
10 works the same as the CD field 404 does with respect to operations of the fields 108 through 110 and the conditional execution field 403.

Registers in the memory unit 3 shown in Figure 1 will be described.

15 The register 30 in the PC controlling part 13, the register 40 in the memory controlling part 14, the register 50 in ALU 15 and the register 60 in the shifter 16 hold a control signal 11 obtained by decoding operation_0 designated by the operation field 106
20 without change, respectively.

The register 31 in the PC controlling part 13, the register 41 in the memory controlling part 14, the register 51 in ALU 15 and the register 61 in the shifter 16 hold a description designating a condition for
25 executing an operation of operation_0 respectively. In this embodiment, a CC value of the condition execution field 401 in the instruction format 101 is held without

change.

The register 32 in the PC controlling part 13, the register 42 in the memory controlling part 14, the register 52 in ALU 15 and the register 62 in the shifter 5 16 hold a description of a time for judging the execution condition of the operation of operation_0. Specifically, the PC value (address value) at a time for judging the execution condition of the operation of operation_0 is held. For example, an address value which 10 should be held in the register 32 is produced by the PC controlling part 13 based on the CD field 404 having an offset value. The PC controlling part 13 receives the offset value of the CD field 404 from the instruction decode unit 2, adds the offset value to the address of 15 instruction format 101 in which operation_0 is described and sets the added value in the register 32. When the CD format 404 designates a register number, the PC controlling part 13 receives the register number from the instruction decode unit and sets it in the register 20 32 directly. In the register 32 specified by the register number held in the register 32 holds a PC value (address value) to be used at the time for judging the execution condition of the operation of operation_0. Incidentally, a bit for distinguishing whether a PC 25 value is held in the register 32 or a register number is held therein is provided in the register 32.

The memory controlling part 14, ALU 15 and the

shifter 16 produce a PC value or a register number in a similar manner to that in the PC controlling part 13 and set these respectively in the register 42, the register 52 and the register 62.

5 The PC controlling part 13 works as an operation unit of the instruction execution unit when operation_0 designated by the operation field 106 is a branch sub-instruction as described in the above (A-9). In a case that operation_0 is a branch sub-instruction, the
10 registers 30 through 32 function as described in the above, and other registers 40 through 42, 50 through 52 and 60 through 62 are not used. Also, the memory controlling part 14 works as an operation unit of the instruction execution unit when operation_0 designated
15 by the operation field 106 is a memory access sub-instruction such as a load/store sub-instruction described in the above (A-1). In a case that operation_0 is such a memory access sub-instruction, the registers 40 through 42 function as described in the above, and
20 other registers are not used.

 ALU 15 works as an operation unit of the instruction execution unit when operation_0 designated by the operation field 106 is an arithmetic operation sub-instruction as described in the above (A-5) or a logical
25 operation sub-instruction as described in the above (A-6). In a case that operation_0 is an arithmetic operation sub-instruction or a logical operation sub-

instruction, the registers 40 through 42 function as described in the above, and other registers are not used.

Further, the shifter 16 works as an operation unit
5 of the instruction execution unit when operation_0
designated by the operation field 106 is a shift
operation instruction as described in the above (A-7).
In a case that operation_0 is such a shift operation
sub-instruction, the registers 40 through 42 function as
10 described in the above, and other registers are not
used.

A register 33 and a memory circuit 34 in the PC
controlling part 13 are used when operation_0 is
specifically a delayed branch sub-instruction and a
15 delayed jump sub-instruction.

Figure 7 is a diagram for explaining a basic format
320 of delayed branch sub-instruction. Basically, the
format 320 of delayed branch sub-instruction includes an
operation code 321, a field for designating amount of
20 execution delay 322 which designates an amount of delay
by which a time of executing branch is delayed and a
field for designating branch destination 323 which
designates an offset or an address for designating a
branch target address. The delayed branch sub-
25 instruction is realized by, for example, the format 116
(Short_D1), the format 117 (Short_D2) or the format 118
(Long). The format 116 (Short_D1) is used when a

register set value is used as the amount of delay. The
format 117 (Short_D2) is used when an immediate value is
used as the amount of delay. The format 118 (Long) is
used when the branch target address is designated by a
5 32-bit immediate value. In these formats, an operation
code is designated by the field 120. Further, the field
129 is used to designate whether the field 128 indicates
a register number or an immediate value. The field 121
is used as a region for designating register when an
10 amount of delay is designated by a register in each sub-
instruction of DBRA, DBSR, DJMP, and DJSR. The field 131
is used as a region of immediate value designating the
amount of delay.

The delayed jump sub-instruction is also described
15 in the format shown in Figure 7. However, the field 323
designates a register number of register holding an
address of jump destination.

In this embodiment, the amount of delay described in
the field 131 as an immediate value is an offset value
20 OVB from an address X of dual operation instruction 101
in which a delayed branch sub-instruction or a delayed
branch sub-instruction is described. Accordingly, the
amount of delay can be variably set when an user
appropriately sets a value in the field for designating
25 amount of execution delay 322. However, it is necessary
for the user to set the field for designating amount of
execution delay 322 or the CD field 404 so that a time

for executing the branch is not earlier than a time for judging the branch condition determined by the CD field 404.

The register 33 provided in the PC controlling part 5 13 holds a description about a time for executing the branch designated by the delayed branch sub-instruction or the delayed jump sub-instruction. Specifically, a PC value (address value) at the time for executing the branch is held. The address value to be held in the 10 register 33 is produced by the PC controlling part 13 based on the field 322 having an offset value. The PC controlling part 13 receives the offset value of field 322 from the instruction decode unit 2, adds the offset value to an address of the instruction format 101 in 15 which a delayed branch sub-instruction or a delayed jump sub-instruction is described and sets the added value to the register 33. When the format 322 designates a register number, the PC controlling part 13 receives the register number from the instruction decode unit and 20 sets it directly to the register 33. The register specified by the register number held in the register 33 holds the PC value at the time of executing the branch.

The memory circuit 34 of the PC controlling part 13 is to hold a description indicating whether or not an 25 execution condition is satisfied as a result of judging the execution condition of the delayed branch sub-instruction or the delayed jump sub-instruction by the

PC controlling part 13.

The control signal 11 held in the register 30 includes a description of 1 bit for distinguishing whether a sub-instruction with respect to the control
5 signal 11 is a delayed branch sub-instruction or an ordinary branch sub-instruction without delay.

Further, registers in the integer operation unit 4 shown in Figure 1 will be described.

Each of the register 70 of the multiplier 17, the
10 register 80 of ALU 19 and the register 90 of the shifter 20 holds the control signal 12 obtained by decoding operation_1 designated by the operation field 106 without change.

Each of the register 71 of the multiplier 17, the
15 register 81 of ALU 19 and the register 91 of the shifter 20 holds a description which designates an execution condition of an operation of operation_1. In this embodiment, a CC value of the condition execution field 402 in the instruction format 101 is held without
20 change.

Each of the register 72 of the multiplier 17, the register 82 of ALU 19 and the register 92 of the shifter 20 holds a description at a time of judging an execution condition of the operation of operation_1. When the CD
25 field 405 has an offset value, it is held as a PC value (address value) to be used at the time of judging the execution condition of the operation of operation_1. The

PC value to be held is a value obtained by adding an address of the instruction format 101 in which operation_1 is described to the offset value. Also, when the CD field 405 designates a register number, the
5 register number is held directly. A register specified by the register number which is held in the register 72 holds a PC value (address value) to be used at the time of judging the execution condition of the operation of operation_1. The multiplier 17, ALU 19 and the shifter
10 20 set a value respectively to the register 72, the register 82 and the register 92 in accordance with the CD field 405 receiving a content from the instruction decode unit.

The multiplier 17 works as an operation unit of the
15 instruction execution unit when operation_1 designated by the operation field 107 is a multiply sub-instruction and a multiply and add sub-instruction, both of which accompany a multiplication as shown in the above (B-1). In a case that operation_1 is a sub-instruction
20 accompanying a multiplication, the registers 70 through 72 function as mentioned in the above, and other registers 80 through 82 and 90 through 92 are not used. ALU 19 works as an operation unit of the instruction execution unit when operation_1 designated by the
25 operation field 107 is an arithmetic operation sub-instruction as in the above (A-5) or a logical operation sub-instruction as in the above (A-6). In the case that

operation_1 is an arithmetic operation sub-instruction or a logical operation instruction, the registers 80 through 82 function as mentioned in the above and other registers are not used.

5 The shifter 20 works as an operation unit of the instruction execution unit when operation_1 designated by the operation field 107 is a shift operation sub-instruction as in the above (A-7). In the case that operation_1 is a shift operation instruction, the
10 registers 40 through 42 function as mentioned in the above, and other registers are not used.

 In a case that a single operation instruction is processed by the instruction format 102, one of the PC controlling part 13, the memory controlling part 14, ALU
15 15, ALU 19, the shifter 16, the shifter 20 and the multiplier 17 executes an operation sub-instruction as an operation unit of the instruction execution unit depending on a type of operation such as a branch, a memory access and an arithmetic operation. The CD field
20 406 designates an amount of delay by which a time for judging an execution condition designated by the execution condition field 403 similarly in a pipeline process of a single operation instruction. Further, a register provided in the operation unit of the
25 instruction execution unit for executing the single operation instruction 102 holds a value having the same content as described about the single operation

instruction 102 in the above.

In the next, operation of the microprocessor 1 will be described in reference of an example of program shown in Figure 8.

5 In this program, a pair of sub-instructions in each row is described by a dual operation instruction 101 having the instruction format shown in Figure 2a, wherein sub-instructions I01, I11, I21, I31, I41, I51, and I61 are described in the operation field 106 as
10 operation_0, and sub-instructions I02, I12, I22, I32, I42, I52, and I62 are described in the operation field 107 as operation_1. Each dual operation instruction is accessed by means of an address number of memory described in the identical row. For example, the sub-
15 instructions I01 and I02 are stored in a memory area of No. 1000 through No. 1007 and accessible at an address No. 1000.

The sub-instruction I01 is a branch instruction BRA for taking a branch to sub-instruction I11 and I12
20 having a description of "loop" when a branch condition that "flag F0 is false (namely, flag F0 holds 0) is satisfied, wherein the branch condition is judged at the time of executing the sub-instruction I41 and I42. And the branch instruction BRA is an ordinary branch
25 instruction without delay. The sub-instruction I21 is an add instruction ADD which adds a content of a register R2 to a content of a register R3 and stores the result

of addition in the register R2. The sub-instruction I31 is a comparison instruction CMPEQ which writes "1" in the flag F0 when the content of register R2 and a content of register R4 are equal and "0" therein when the contents are not equal. The sub-instructions I11, I41, I51, and I61 are arbitrary arithmetic operation instructions by which the stage E/M is processed in the memory unit 2, and other sub-instructions are arbitrary arithmetic operation instructions by which the stage E/M is processed in the integer operation unit 3.

On the execution condition field 401 corresponding to the sub-instruction I01, there is described "CC=010"; and sub-instructions other than the sub-instruction I01 are instructions executed unconditionally, wherein execution condition fields corresponding thereto have a description of "CC=000". Further, in the CD field 404 corresponding to the sub-instruction I01, an offset value '20' is described. Further, dual operation instructions described in the rows are instructions which execute two operations of the sub-instructions in parallel by setting the FM fields 103 and 104 to be 00.

In the program shown in Figure 8, one loop is formed by the four dual operation instructions in the address Nos. 1008, 1010, 1018, and 1020, wherein this program means that the four dual operation instructions are repeatedly executed in a sequential manner until the flag F0 becomes true when a comparison instruction is

executed by the sub-instruction I31.

Figure 9 shows operation of the microprocessor 1 which processes the program shown in Figure 8 in pipeline. In the Figure, clocks t1 through t13 designate continuous one clock cycles, and all pipeline stages in each clock cycle are processes in parallel. For example, in clock t5, each stage W in the sub-instructions I11 and I12, each stage E/M in the sub-instructions I21 and I22, each stage D/A in the sub-instructions I31 and I32 and each stage IF in the sub-instructions I41 and dI42 are processed in parallel respectively. Stages in other clock cycles are similarly processes thereto.

In Figure 9, an address value held by the PC of microprocessor 1 shows that the clock cycle corresponding to the address value is a cycle just after the cycle of processing the dual operation instruction which is accessed by the address value on the instruction decode stage D/A.

Concerning the sub-instructions I01 and I02, the instruction fetch stage IF, and the instruction decode stage D/A are processed in parallel respectively in the clock t1 and the clock t2. Although the instruction execution stage E/M and the write back stage W of the sub-instruction I02 are processed respectively in the clocks t3 and t4, the instruction execution stage E/M of the sub-instruction I01 is not processed by judging the execution condition and branching based on this

judgement, until it is enabled to process.

The instruction decode unit 2 detects that the sub-instruction I01 is a sub-instruction for delaying judgement of execution condition in accordance with the field for designating condition 401 and the CD field 404 both of which are of the branch sub-instruction BRA as the sub-instruction I01, and the contents of the field for designating condition and the CD field are outputted to the PC controlling part as a control signal for judging the execution condition with a delay. Also, the field for designating operation is decoded in the decoder 8 and a control signal 11 is outputted in response to the result of decoding. In the control signal 11, a first description for controlling the PC controlling part 13 so as to take a branch of the branch sub-instruction BRA, a second description for showing that the sub-instruction I01 is an ordinary sub-instruction which does not cause a delay of execution, a third description for indicating a branch address of the branch sub-instruction BRA are included. The third description is the branch address itself, which is calculated by an adder for exclusively calculating address (not shown) based on an offset designated by the field 323 of the branch sub-instruction BRA at the instruction decode stage E/A.

In the clock t3, the PC controlling part 13 receives a description that the sub-instruction I01 from the

instruction decode unit 2 is a conditional sub-
instruction and a sub-instruction for delaying a time of
judging the condition, and holds the control signal 11
with respect to the branch sub-instruction BRA in its
5 register 30 without change. At this time, the branch is
not executed based on the control signal 11. A value of
CC=010 which is the execution condition field outputted
from the instruction decode unit 2 is held in the
register 31 without change. Further, the PC controlling
10 part 13 receives an offset value "20" which is the CD
field from the instruction decode unit 2 and the address
No. 1000 from the PC, adds these, and holds a result of
the addition of the address No. 1020 in its register 32
at the clock t3. The PC controlling part 13 is comparing
15 the value held in the register 32 with the value
indicated by the PC. The PC controlling part 13 judges
the execution condition of the branch sub-instruction
BRA based on a clock cycle at which the address value in
the PC is equal to the address value in the register 32,
20 namely, a CC value held in the register 31 at the clock
t7.

On the other hand, the decoder 9 of the instruction
decode unit 2 analyzes the field for designating
operation 107 to thereby output a control signal 12 for
25 commanding ALU 18 to perform an arithmetic operation,
with respect to the sub-instruction I02. The instruction
decode unit 2 detects, based on the field for

designating condition, that the sub-instruction I02 is an unconditional sub-instruction, and outputs a description of showing that the sub-instruction I02 is unconditional (CC=000).

5 ALU 18 performs an add operation in accordance with a control signal 12 without holding the control signal 12 in its register 80 when the description that it is unconditional is received. Further, the value of execution condition field 402 and the value of CD field 10 405 both in the sub-instruction I01 are outputted from the instruction decode unit 2. However, ALU 18 holds values already held in the registers 81 and 82 without change by ignoring the value of execution condition field 402 and the value of CD field 405. As for other 15 sub-instructions to be executed unconditionally described in the below, the similar processes thereto are applicable.

With respect to the sub-instructions I11 and I12, the instruction fetch stages IF, the instruction decode 20 stages D/A, the instruction execution stages E/M and the write back stage W are processed in parallel respectively at the clocks t2, t3, t4, and t5.

With respect to the sub-instructions I21 and I22, the instruction fetch stages IF, the instruction decode 25 stages D/A, the instruction execution stages E/M and the write back stages W are processed in parallel respectively at the clocks t3, t4, t5, and t6. On the

add sub-instruction ADD which is the sub-instruction I21, a content of the register R5 and a content of the register R6 are operated to add in the stage E/M of the clock t5 by ALU 15 and a result of the addition is
5 written in the register R5 in the stage W of the clock t6.

With respect to the sub-instructions I31 and I32, the instruction fetch stages IF, the instruction decode stages D/A and the instruction execution stages E/M are
10 processed in parallel respectively at the clocks t4, t5, and t6. On the comparison sub-instruction CMPEQ which is the sub-instruction I31, a content of the register R2 and a content of the register R4 are compared by ALU 15 in the stage E/M of the clock t6, wherein if the
15 contents are equal '1' is written in the flag F0 and if not, '0' is written in the flag F0. Although the execution stage E/M of the comparison sub-instruction CMPEQ can not basically be started before the clock t6 at which a result of operation by the add sub-
20 instruction ADD is written in the register R5, it is processed using a result of the operation by the add sub-instruction which is obtained at the clock t5 by a bypass circuit provided in the processor 1.

Incidentally, in the case of the comparison sub-
25 instruction CMPEQ, the write back stage W is not included. On the other hand, the write back stage W of the sub-instruction I32 is processed at the clock t7.

With respect to the sub-instructions I41 and I42, the instruction fetch stages IF, the instruction decode stages D/A, the instruction execution stages E/M and the write back stages W are processed in parallel
5 respectively at the clocks t5, t6, t7, and t8.

The branch condition of the branch sub-instruction BRA is judged at the clock t7 because a result of operation by the sub-instruction I31 should be referred to. At the clock t7, the PC controlling part 13 refers
10 to the flag F0 in accordance with the value of '010' held in the register 31, determines to take a branch when the flag F0 is '0', and determines not to take a branch when the flag F0 is '1'.

The PC controlling part 13 ignores a content of the
15 register 33 in accordance with an event that the second description of the control signal 11 held in the register 30 shows that the branch sub-instruction BRA is a branch sub-instruction without delay. In other words, the instruction execution stage E/M is processed at the
20 same clock t7 in accordance with an event that the branch condition is determined in the PC controlling part 13. Since the sub-instruction I41 is an arithmetic operation sub-instruction and the instruction execution stage E/M therefore is processed in ALU 15 not in the PC
25 controlling part 13, the stages E/M of the sub-instruction I01 and the sub-instruction I41 can be processed in parallel.

In the operation of microprocessor 1 shown in Figure 9, provided that '0' is written in the flag F0 in the comparison sub-instruction CMPEQ, a judgement of condition is completed when the PC controlling part 13 produces a description (for example, a logical value of '1') for showing that the execution condition is satisfied, and a branch is executed based on the control signal 11 held in the register 30 in response to the description. Specifically, the PC controlling part 13 outputs a branch address (in this case, the address No. 1008) to the instruction RAM 6 through the IA bus, and sets the branch address in the PC at the next clock cycle in accordance with the third description of the control signal 11. Further, in accordance with the first description of the control signal 11, the PC controlling part 13 controls the instruction RAM 6 so as to give a dual operation instruction (sub-instructions I11 and I21) stored in the branch address No. 1008 to the instruction decode unit 2. The PC controlling part 13 controls the instruction decode unit 2, the memory unit 3 and the integer operation unit 4 so as to cancel pipeline processes of sub-instructions I51 and I52, which are already processed by the stages IF and D/A at the clock t7, and pipeline processes of sub-instructions I61 and I62, which are already processed by the stage IF, in accordance with the first description. However, the write back stages W of the sub-instructions I41 and

I42, which are processed by the stage E/M at the clock t7, are proceeded to process without canceling.

The instruction decode unit 2 receives the dual operation instruction of the sub-instructions I11 and I12 existing in the address No. 1008, which are outputted from the instruction RAM 6, whereby the instruction fetch stage IF is processed at the clock t8. With respect to the sub-instructions I11 and I12, the instruction decode stages D/A, the instruction execution stages E/M and the write back stages W are processed respectively at the clocks t9, t10 and t11. The PC holds the address No. 1008 without change until the clock 10 at which the instruction execution stage E/M of the dual operation instruction in the address No. 1008 is processed and counts by every number of 8 at and after the clock t10.

Following the dual operation instruction of the sub-instructions I11 and I12, each dual operation instruction in the address No. 1010, 1018, 1020 is processed in the pipeline sequentially interposing an one-clock delay. Since the registers 30, 31, and 32 of the PC controlling part 13 holds the value held at the clock t3 without change, the PC controlling part 13 executes branching of the branch sub-instruction BRA by judging the same execution condition that a value of the flag F0 is '0' in reference of the flag F0 based on the contents of the registers 30, 31, and 32 at the clock

t13 at which the PC holds the address No. 1020 again. Although it is not shown, the execution condition is determined by a renewal of the flag F0 which is a result of executing the sub-instruction I31 obtained at the
5 clock t12.

When it is judged that the execution condition of the branch sub-instruction BRA is not satisfied at the clock t7, the PC controlling part 13 produces a description of showing that the execution condition is
10 not satisfied (for example, a logical value of '0') and a branch is not executed without referring to the control signal 11 held in the register 30 in response to the description. The dual operation instruction of the sub-instructions I51 and I52 and the dual operation
15 instruction of the sub-instructions I61 and I62 are successively processed without canceling the processes of pipeline stage thereof. Further, the contents held in the registers 30, 31, and 32 can be left without change until they are renewed by a next conditional branch sub-
20 instruction, or can be reset completely.

Meanwhile, when the offset value of the CD field 404 with respect to the branch sub-instruction BRA is 0, the execution condition is judged at the clock t3 at which the PC value indicates the address No. 1000 which is the
25 address of the branch sub-instruction itself and the branch is executed.

Another operation of the microprocessor 1 will be

described using an example of program shown in Figure 10.

As in Figure 8, two sub-instructions in each row are described by a dual operation instruction 101 having the instruction format shown in Figure 2a; sub-instructions I01, I11, I21, I31, I41, I51, and I61 are described in the operation field 106 as operation_0; and sub-instructions I02, I12, I22, I32, I42, I52, and I62 are described in the operation field 107 as operation_1.

10 The sub-instruction I01 is a jump sub-instruction DJMP which jumps to an instruction in an address number held in the register R5 when an execution condition that the flag F0 is true (holding '1') is satisfied, wherein the jump instruction DJMP is a delayed jump sub-

15 instruction having a format shown in Figure 7. The sub-instruction I11 is an add sub-instruction ADD which writes a result of adding a content of the register R1 to a content of the register R2 in the register R1. The sub-instruction I21 is a comparison sub-instruction

20 CMPEQ which compares a content of the register R1 with a content of the register R3, writes '1' in the flag F0 if these are equal, and writes '0' in the flag F0 if these are not equal. The sub-instruction I31 is an add sub-instruction which adds a content of the register R5 to a

25 content of the register R6 and writes a result of the addition in the register R5. The sub-instructions I41, I51, I61, and I71 are arbitrary arithmetic operation

sub-instructions of which stages E/M are processed in the memory unit 3, and other sub-instructions are arbitrary arithmetic operation sub-instructions of which stages E/M are processed in the integer operation unit

5 4.

In the execution condition field 401 corresponding to the sub-instruction I01, "CC=001" is described. And sub-instructions other than the sub-instruction I01 are sub-instructions which are executed unconditionally.

10 Also, in the CD field 404 corresponding to the sub-instruction I01, an offset value "18" is described.

Further, in a field for designating amount of execution delay (Figure 7) of the sub-instruction I01, an offset value '28' is described as an immediate value.

15 Further, the dual operation instruction described in every row is an instruction for executing operations of two sub-instructions in parallel by setting "00" in the FM fields 103 and 104.

In the program of Figure 10, a jump executed by the jump sub-instruction DJMP of the sub-instruction I01 is executed by judging the value of flag F0, which is a result of comparison by the sub-instruction I21, and an address of jump destination is determined by a result of addition in the sub-instruction I31. The dual operation

20 instructions in the address Nos. 1020 and 1028 are executed despite whether or not a jump is taken by the jump instruction DJMP.

25

Figure 11 shows operation of the microprocessor 1 which processes the program shown in Figure 10 by a pipeline processing. As shown in Figure 9, the clocks t1 through t13 respectively designate sequential one clock
5 cycles, wherein processings of all pipeline stages in each clock cycle are conducted in parallel. The PC value has the same meaning as that in the Figure 9.

With respect to the sub-instructions I01 and I02, the instruction fetch stages IF and the instruction
10 decode stages D/A are processed in parallel respectively at the clocks t1 and t2. Although the instruction execution stage E/M and the write back stage W of the sub-instruction I02 is processed respectively at the clocks t3 and t4, the instruction execution stage E/M of
15 the sub-instruction I01 is not subjected to a judgement of execution condition and not executed to branch based on the judgement before a permission is obtained.

The delayed jump sub-instruction DJMP as the sub-instruction I01 is decoded by the decoder 8 shown in
20 Figure 1, whereby a control signal 11 is outputted in response to a result of this decoding. In this control signal 11, a first description for controlling the PC controlling part 13 so as to take a branch of a jump sub-instruction DJMP, a second description for showing
25 that the sub-instruction I01 is a sub-instruction of causing a delay of execution and a third description for showing a branch target address of the branch sub-

instruction BRA are included. The third description is a register number designated by the jump sub-instruction DJMP.

At the clock t3, the PC controlling part 13 holds
5 the control signal 11 to the jump sub-instruction DJMP
in the register 30 without change and does not execute a
branch based on the control signal 11. A value of
CC=001, which is the execution condition field of the
jump sub-instruction DJMP outputted from the instruction
10 decode unit 2, is held in the register 31 without
change. Further, the PC controlling part 13 receives an
offset value "18" of the CD field of the jump sub-
instruction DJMP from the instruction decode unit 2 and
the address No. 1000 from the PC respectively at the
15 clock t3 and adds these, wherein the register 32 holds a
result of the addition, namely the address No. 1018. The
PC controlling part 13 is comparing a value held in the
register 32 and a value indicated by the PC. The PC
controlling part 13 judges the execution condition of
20 the jump sub-instruction DJMP at the first time based on
the CC value held in the register 31 at a clock cycle at
which the address value of the PC is equal to the
address value of the register 32, namely at the clock
t6.
25 Further, the PC controlling part 13 receives an
offset value "28", which is the field for designating
amount of execution delay 322 of the jump sub-

instruction DJMP, from the instruction decode unit 2,
receives the address No. 1000 from the PC at the clock
t3 and adds these, wherein the register 33 holds a
result of the addition, namely the address No. 1028. The
5 PC controlling part 13 is comparing a value held in the
register 33 with a value indicated by the PC. A jump by
the jump sub-instruction DJMP is executed at the first
time at a clock cycle at which the address value of the
PC is equal to the address value of the register 33,
10 namely the clock t8.

With respect to the sub-instructions I11 and I12,
the instruction fetch stages IF, the instruction decode
stages D/A, the instruction execution stages E/M and the
write back stages W are processed in parallel
15 respectively at the clocks t2, t3, t4, and t5. In the
add sub-instruction ADD of the sub-instruction I11, a
content of the register R1 and a content of register R2
are added by ALU 15 in the stage E/M, and a result of
the addition is written in the register R1 in the stage
20 W.

With respect to the sub-instructions I21 and I22,
the instruction fetch stages IF, the instruction decode
stages D/A and the instruction execution stages E/M are
processed in parallel respectively at the clocks t3, t4
25 and t5. In the comparison sub-instruction CMPEQ as the
sub-instruction I21, a content of the register R1 and a
content of the register R3 are compared by ALU 15 in the

execution stage E/M of the clock t5; '1' is written in the flag F0 if these are equal; and '0' is written in the flag F0 if these are not equal.

In a case of the comparison sub-instruction CMPEQ,
5 the write back stage W does not exist. On the other hand, the write back stage W of the sub-instruction I32 is processed at the clock t7.

The execution condition of the jump sub-instruction DJMP is judged at the clock t6 because the flag F0 which
10 is the result of operation of the sub-instruction I21 should be referred to. At the clock t6, the PC controlling part 13 refers to the flag F0 in accordance with a value of '001' held in the register 31,
determines to take a branch if the flag F0 is '1', and
15 determines not to take the branch if the flag F0 is '0'. At the clock t6, a description of one bit for determining whether or not the branch is taken in the memory circuit 34 is set at the clock t6. However, the actual jump by the jump sub-instruction is not executed
20 before the clock t8.

In this, a case that '1' is written in the flag F0 by the comparison sub-instruction CMPEQ and thereby the execution condition is satisfied will be considered.

In the memory circuit 34, '1' which is the
25 description that the execution condition is satisfied is set. In accordance with the second description for showing that the jump sub-instruction DJMP of the

control signal 11 held in the register 30 is a delayed jump sub-instruction, the PC controlling part 13 checks whether or not the address value held in the register 33 is still in agreement with the PC value. If it is
5 determined that the execution condition is satisfied, since these are not in agreement with at the stage of clock t6, the PC controlling part 13 does not execute the branch in accordance with the first description and the third description of the control signal 11 held in
10 the register 30.

With respect to the sub-instructions I31 and I32, the instruction fetch stages IF, the instruction decode stages D/A, the instruction execution stages E/M and the write back stages W are processed in parallel
15 respectively at the clocks t4, t5, t6, and t7. In the add sub-instruction DD of the sub-instruction I31, a content of the register R5 and a content of the register R6 are added in the stage E/M and a result of the addition is written in the register R5 in the stage W.

20 With respect to the sub-instructions I41 and I42, the instruction fetch stages IF, the instruction decode stages D/A, the instruction execution stages E/M and the write back stages W are processed in parallel respectively at the clocks t5, t6, t7, and t8. With
25 respect to the sub-instructions I51 and I52, the instruction fetch stages IF, the instruction decode stages D/A, the instruction execution stages E/M and the

write back stages W are processed in parallel respectively at the clocks t6, t7, t8, and t9.

The jump by the jump sub-instruction DJMP is executed at the clock t8 since the register R5 which
5 indicates a result of the operation by the sub-instruction I31 should be referred to. At the clock t8 held in the register 33 is in agreement with PC value. Then, in accordance with an even that '1' is held in the memory circuit 34, the PC controlling part 13 processes
10 the instruction stage E/M of the jump sub-instruction DJMP at the clock t8. The PC controlling part 13 takes a branch based on the control signal 11 held in the register 30. Since the sub-instruction I51 is an arithmetic sub-instruction and the processing of the
15 instruction execution stage E/M is performed in ALU 15 not in the PC controlling part 13, the stages E/M of the sub-instructions I01 and I51 can be processed in parallel.

Specifically, in accordance with the third
20 description of the control signal 11, the PC controlling part 13 outputs the branch target address (for example, the address No. 2000) held in a register designated by the third description to the instruction RAM 6 through the IA bus and sets the branch address in the PC at the
25 next clock cycle. Further, in accordance with the first description of the control signal 11, the PC controlling part 13 controls the instruction RAM 6 so as to give the

dual operation instruction stored in the branch target address No. 2000 to the instruction decode unit 2.

Further, the PC controlling part 13 controls the instruction decode unit 2, the memory unit 3 and the integer operation unit 4 so as to cancel the pipeline processings of the sub-instructions I61 and I62, of which the stages IF and D/A are already processed and the sub-instructions I71 and I72, of which stages IF are processed at the clock t8 in accordance with the first description. However, the write back stages W of the sub-instruction I51 and I52, of which stage E/M are processes at the clock t8, are processed without cancellation.

The instruction decode unit 2 receives the dual operation instruction in the address No. 2000 outputted from the instruction RAM 6, wherein each instruction fetch stage IF of the dual operation instruction is processed at the clock t9, and succeedingly, the stages D/A, E/M and W are sequentially processed at each one clock cycle.

Further, in a case that the execution condition of the jump sub-instruction is not satisfied at the clock t6, a description of '0' which means that the execution condition of the jump sub-instruction is not satisfied is held in the memory circuit 34 at the clock t6. In addition, if the address value held in the register 33 at the clock t8 is equal to the PC value, the PC

controlling part 13 does not take a branch without referring to the control signal 11 held in the register 30 in accordance with an event that '0' is held in the memory circuit 34. Processing in each pipeline stage of the dual operation instruction of the sub-instructions I61 and I62 and the dual operation instruction of the sub-instructions I71 and I72 is successively conducted without cancellation. Contents held in the registers 30, 31, 32, and 33 and the memory circuit 34 can be left as these are, until they are renewed by a next conditional branch sub-instruction, or can be reset completely.

As another method, it is possible to rewrite from the control signal 11 held in the register 30 to a control signal by which the PC controlling part 13 is controlled to invalidate the jump, namely, not to execute the jump, in accordance with a description indicating that the condition is not satisfied by judging the condition at the clock t6. If the description showing that the condition is satisfied is produced, the control signal 11 of the register 30 is held without change. In this case, it is not necessary to provide the memory circuit 34 which is referred to at the time of executing the jump actually at the clock t8.

Needless to say, as for the delayed jump sub-instruction DJMP or the delayed branch sub-instruction DBRA, it is possible to execute a branch at the same cycle as the clock cycle at which the execution

condition is judged by changing a value of the field for designating amount of execution delay 322.

It is not limited to a conditional branch or a jump sub-instruction to delay a time for judging an execution condition. It is possible to delay a time for judging execution conditions designated by the fields for designating condition 401 through 403 by the CD formats 404 through 406 in execution processings of two arbitrary sub-instructions of a dual operation instruction and of an arbitrary single operation instruction 102. This can be realized by providing the registers 40 through 42 of the memory controlling part 14, the registers 50 through 52 of ALU 15, the registers 60 through 62 of the shifter 16, the registers 70 through 72 of the multiplier 17, the registers 80 through 82 of ALU 19 and the registers 90 through 92 of the shifter 20 every registers having the same structure and the same function as those of the registers 30 through 32 of the PC controlling part 13.

For example, a case that the sub-instruction I12 of the program shown in Figure 8 is represented by the following formula A will be considered.

FORMULA A

I12(ADD F1T R13 R8, R8, R9)

25

This is an add sub-instruction which adds a content of the register R8 to a content of the register R9 when

an execution condition that the flag F1 is true is satisfied and writes a result of the addition in the register R8, wherein a judgement of the execution condition is conducted at a clock cycle at which the PC value is the address value held in the register R13. CC=011 is described in the field for designating condition 402 and the register No. '13' is described in the CD field 405. The address No. 1018 is held in the register 13.

10 In Figure 9, the instruction execution stage E/M of the sub-instruction I12 is processed at the clock t4. However, if the sub-instruction I12 is an add sub-instruction as shown in Formula A, the instruction execution stage E/M is processed at the clock t6.

15 The sub-instruction I12 is decoded by the decoder 9 at the clock t3, wherein the control signal 12 is outputted in response to a result of the decoding. In this control signal 12, there is included a first description for controlling ALU 19 so as to execute an add sub-instruction and a second description for designating the registers R8 and R9 which are used for the operation.

20 At the clock t4, ALU 19 holds the control signal 12 in its register 80 without change and does not execute the addition. A value of CC=011 which is the execution condition field outputted from the instruction decode unit 2 is held in the register 81 without change.

Further, ALU 19 receives the register No. '13' which is the CD field from the instruction decode unit 2 and holds it without change. ALU 19 is comparing a value held in the register of the number '13' held in the
5 register 82 with a value indicated by the PC. ALU 19 judges an execution condition of the add sub-instruction at the first time based on the CC value held in the register 81 at the clock t6, which is the clock cycle at which an address value in the PC is equal to the address
10 No. 1018. ALU 19 produces a description showing whether or not the execution condition is satisfied.

ALU 19 performs the add operation at the same clock t6 based on the control signal 12 held in the register 80 according to the description and makes a designated
15 general purpose register 5 hold a result of the addition at the clock t7 when the condition is satisfied. When the condition is not satisfied, the control signal 12 is ignored and the add operation is not performed.

As another means applicable when the condition is
20 not satisfied, although ALU 19 can perform the add operation based on the control signal 12, ALU 19 is constituted so as not to write the result of the addition in the register designated by the add operation instruction in accordance with the description showing
25 whether or not the execution condition is satisfied.

An operation sub-instruction, namely, a comparison sub-instruction which compares two values and reflects a

result of whether these are equal or non-equal or a
result of which is larger in the flag F1, for
determining a value of the flag F1, which is the
execution condition of the add sub-instruction of
5 Formula A, can be put in a sub-instruction positioned in
a lower address than the sub-instruction I12 in the
program sequence, for example, the position of the sub-
instruction I22.

In each of the CD fields 404, 405, and 406,
10 descriptions showing that execution conditions
designated by fields for designating execution condition
are judged in the instruction decode stage D/A of
operation sub-instructions described in each operation
field can be described in addition to the description of
15 offset value and the description which designate
registers.

For example, in Figure 8, when the CD field 404
corresponding to the branch sub-instruction BRA of the
sub-instruction I01 has a description indicating that an
20 execution condition is judged in the instruction decode
stage D/A, the instruction decode unit 2 starts the
judgement of the execution condition in response to the
CD field 404 and refers to a content of the designated
flag F0 in PSW10 in accordance with the field for
25 designating condition 401 at the clock t2. Then the
instruction decode unit 2 outputs a fourth description
showing whether or not the execution condition of the

sub-instruction I01 is satisfied. Further, the instruction decode unit 2 forms a control signal 11 based on a result of decoding by the field for designating operation of the branch sub-instruction BRA by the decoder 8. In the control signal 11, there are included a first description for controlling the PC controlling part 13 so as to take a branch of the branch sub-instruction BRA, a second description representing an ordinary instruction which does not cause a delay of branch execution and a third description indicating a branch target address of the branch sub-instruction BRA.

In a case that the execution condition is satisfied, the instruction decode unit 2 outputs a description that the execution condition is satisfied as the fourth description, namely, it can be executed unconditionally in the PC controlling part 13. The PC controlling part 13 receives the description indicative of the unconditionalness and conducts the branch in accordance with the control signal 11 at the clock t3 without holding the control signal 11 in the register 30.

In a case that it is judged that the execution condition is not satisfied, the instruction decode unit 2 outputs a description indicating that the operation is invalid as the fourth description. The PC controlling part 13 does not execute the branch even if the control signal 11 is received in accordance with the description indicating that the operation is invalid.

As for other kinds of operating sub-instructions than the branch sub-instruction, the above process is similarly applicable. However, in an arithmetic operation such as an addition, in the case that the
5 operation is judged to be invalid in the instruction decode stage D/A, it is possible not to write a result of the operation in a register designated by the operation instruction even though the operation unit executes the arithmetic operation.

10 In addition, if instructions such as the branch sub-instruction DBRA and the jump sub-instruction DJMP, both of which having a function of delay, are not prepared in the microprocessor shown in Figure 1, it is not necessary to provide the register 33 and the memory
15 circuit 34 in the PC controlling part 13.

Although the registers as an means for delaying a time for judging execution conditions are provided respectively in the PC controlling part 13, the memory controlling part 14, ALU 15, ALU 19, the shifter 16, the
20 shifter 20, and the multiplier 17, the registers can be provided in a selected part of these operation units of the instruction execution unit.

For example, among the operation sub-instructions to be executed under a predetermined condition, a
25 conditional branch sub-instruction or a conditional jump sub-instruction is used most frequently in a program made by a user. Accordingly, in a processing of only a

branch sub-instruction, a function of delaying a time for judging the execution condition may be used. In this case, the registers 30, 31, and 32 are provided in only the PC controlling part 13.

5 In the above, when the execution condition field designates the unconditional execution of CC=000, the value of the CD field corresponding to the execution condition field is ignored by the unit 2, 3 and 4 and the sub-instruction corresponding to the execution
10 condition field is decoded and executed in continuous one clock cycles. However, without disregarding the CD field, the sub-instruction may be decoded and unconditionally executed by delay the value of the CD field after decoding.

15 As mentioned in the above, this embodiment has the following characteristics.

(1) For example, in a case of an operation sub-instruction to be executed unconditionally such as the sub-instruction I31 shown in Figure 8 or an operation
20 sub-instruction of which condition is determined in its decode stage even though it is conditional, the instruction is decoded in the period of clock t5; and an operation designated by the sub-instruction I31 is
25 executed in the period of clock t6 succeeding to the period of clock t5, within a delay amount of the above conditional operation sub-instruction. On the other hand, the conditional operation sub-instruction such as

the sub-instruction I01 is decoded in the period of clock t2; and a judgement of the execution condition of branch operation is started at the clock t7 after a period at least longer than the clock t6, which is a period of processing the instruction execution stage E/M such as the sub-instruction I31, from the clock t2. In other words, a time for judging the execution condition is delayed by a delay amount of the period between the clock t3 and the clock t6. Accordingly, as in the program sequence shown in Figure 8, the conditional operation sub-instruction can be put in a point earlier than the operation sub-instruction (the sub-instruction I31), which performs the operation for determining the execution condition of the operation sub-instruction. In a conventional processor, there is a possibility that a sub-instruction other than the NOP sub-instruction cannot be put in the position of the sub-instruction I01. However, when the processor described in this embodiment is adopted, a conditional operation sub-instruction can be put as the sub-instruction I01 and a scheduling of instructions can be flexible.

(2) For example, it is possible to variably set a delay amount representing a time for judging execution conditions of operations held in the registers 32, 42, and so on of each operation unit. Accordingly, for example in the sub-instruction I01, the execution condition is judged at various clock cycles including

the clock t4. Therefore, it is possible to appropriately change the position of conditional operation sub-instruction in accordance with a content of the program at the time of scheduling the sub-instructions. In particular, because values held in the registers 32, 43, and so on are set in accordance with contents described in the CD fields 403, 404, and 405 of a sub-instruction format constituting a program, a programmer or a compiler can easily determine an extent of the delay amount in response to the conditional operation sub-instruction.

(3) Various formats can be used for a description as the delay amount held in the registers 32, 42, and so on. It is possible to hold a clock number as the delay amount. In this case, the registers 32, 42, and so on are constituted as a subtraction counter. For example in the PC controlling part 13, the clock number of 4 is held in the register 32 at the clock t3; the clock number is subtracted in accordance with the counter; and starts a judgement of condition at the clock t7 at which the clock number of 0 is held. However, in this embodiment, an address value is held as the delay amount representing the time for judging the execution condition of operation; and the execution condition is judged at a clock cycle equal to a value of the program counter of the address value. Accordingly, it is more effective because the delay amount is controlled by the

PC value. For example, in a case that the program shown in Figure 8 is constituted to jump to another instruction between the address No. 1008 and the address No. 1020 and return thereafter, it is not necessary to
5 adjust a time for judging execution by changing the values of the register values 32, 42, and so on, whereby a control becomes easy. If the above clock number is held, it is necessary to apply measures such that a subtracted value at the time of jumping to another
10 instruction is evacuated.

(4) In a case of a conditional delayed branch sub-instruction or a conditional delayed jump sub-instruction, a branch condition is judged at the clock t6 as shown in Figure 11 and the branch is executed at
15 the clock t8 which starts after elapsing one clock cycle or more. If the sub-instructions I41 and I42 are sub-instructions which rewrites the flag 0 which determines the execution condition of the jump sub-instruction DJMP at the clock t7, it is possible to use a jump sub-
20 instruction without delay, which executes the jump at the same clock cycle as that for judging the execution condition, as the sub-instruction I01. Accordingly, the conditional delayed operation sub-instruction as described in this embodiment can make an instruction
25 scheduling in a program more flexible.

(5) Each operation unit, for example the PC controlling part 13, holds the control signal 11 for controlling the

PC controlling part 13 so as to execute a first operation sub-instruction, a value of field for designating condition as the first description which indicates the execution condition of the first operation sub-instruction and an address value or a description designating a register holding the address value as the second description which indicates a time for judging the execution condition respectively in its registers 30, 31, and 32. Further, these register values are continuously held without change until they are rewritten by the second conditional operation sub-instruction which is decoded by the instruction decoder 2 after the first operation sub-instruction, in order to delay the time for judging the condition. This procedure is effective when a loop process formed by the conditional operation sub-instruction is conducted. If a description is temporarily held in the registers 30, 31, and 32 beforehand, the conditional operation sub-instruction is executed by using only the descriptions held in the registers 30, 31, and 32 at each occurrence of loop. According to conventional techniques, it was necessary to decode by a decoder after fetching a conditional operation sub-instruction from an instruction RAM at each occurrence of loop. However, in this embodiment, it is sufficient to fetch the operation sub-instruction of the sub-instruction I01 only once and decode the same, whereby the processing efficiency of

microprocessor can be high.

- (6) In this embodiment, all sub-instructions to be processed by the microprocessor has the format shown in Figure 2; fields for designating condition are provided
5 corresponding to fields for designating operation therein; and further fields for designating amount of delay for judging condition are provided corresponding to the fields for designating condition. The fields for designating condition can designate an unconditional
10 execution. In other words, because a space for inserting a description designating whether the sub-instruction is conditional or not, a content of the condition if conditional and a description for indicating a time for judging the condition is secured, whereby a programmer
15 or a compiler can easily constitute sub-instructions.
- (7) In the fields for designating amount of delay for judging condition, it is possible to designate judging in the decode stage of sub-instruction by the instruction decode unit 2 as a time for judging the
20 condition. The instruction execution unit performs an operation as a sub-instruction executable unconditionally if the condition of the sub-instruction is satisfied. This also gives flexibility to a scheduling of the instruction.
- (8) Further, the following modification can be
25 considered. The value referred to at the time of determining the condition in processing the conditional

operation sub-instruction is the flag F0. However, it is not limited to this and the referred value can be a register value holding a plurality of bits.

The microprocessor in this embodiment processes a plurality of pipelines in parallel by adopting a VLIW architecture. However, a function of delaying the time for judging condition in this conditional operation sub-instruction is effective if it is provided in the processor which processes a single pipeline.

(9) Each operation unit is provided with a set of three registers which are a means for delaying the time for judging condition. However, it is possible to further provide one set or plurality sets of three registers having the same function as these three registers in each operation unit. This structure is effective in a case that one loop or plurality loops exist in the loop formed by the program shown in Figure 8.

A program constituted as follows will be considered as an example:

Address No. 100: I01 (BRA F0F #H30 110)
Address No. 108: I11 (BRA F1T #H18 118)
Address No. 110: I21
Address No. 118: I31
Address No. 120: I41
Address No. 128: I51
Address No. 130: I61

In this, the sub-instruction I01 is a branch sub-

instruction which branches to the address No. 110 if the
flag F0 is false, wherein the branch is executed by
judging an execution condition at a cycle of instruction
execution stage of the sub-instruction I61. The sub-
5 instruction I11 is a branch sub-instruction which
branches to the address No. 118 if the flag F1 is true,
wherein the branch is executed by judging an execution
condition in an instruction execution stage of the sub-
instruction I41 in the address No. 120. A first loop is
10 formed by the address Nos. 110 through 130 and a second
loop is formed by the address Nos. 118 through 120 in
this program.

In a case that the sub-instruction I01 is decoded, a
control signal 11 with respect to the branch sub-
15 instruction of the sub-instruction I01, a description
designating the condition and a description designating
a time of judging the condition are held in the
registers 30 through 32 of the PC controlling part 13.
In a case that the sub-instruction I11 is decoded
20 succeeding to the sub-instruction I01, the three
registers which are further provided in the PC
controlling part 13 hold a control signal 11 with
respect to the branch sub-instruction of the sub-
instruction I11, a description designating the condition
25 and a description designating a time for judging the
condition. By this, the two conditional operation sub-
instructions can be executed appropriately by referring

to only the descriptions held in the registers 30, 31, and 32 and the descriptions held in the three registers having the same function respectively, at each occurrence of the first loop and the second loop.

5 Therefore, it is not necessary to fetch the sub-instructions I01 and I02 at every occurrence of the loops and to decode the sub-instructions, whereby a processing performance of the processor can be enhanced.

EMBODIMENT 2

10 Figure 12 shows a constitution of a microprocessor other than that shown in Figure 1.

In this microprocessor, a register 43 and a memory circuit 44 are provided in a memory controlling part 14; a register 53 and a memory circuit 54 are provided in
15 ALU 15; a register 63 and a memory circuit 64 are provided in a shifter 16; a register 73 and a memory circuit 74 are provided in a multiplier 17; a register 83 and a memory circuit 84 are provided in ALU 19; and a register 93 and a memory circuit 94 are provided in a
20 shifter 20. A part of the constitution which has not been described in the above is the same as that shown in Figure 1.

In the microprocessor shown in Figure 1, the operation sub-instruction having a delay function,
25 namely an operation sub-instruction which executes the operation at a clock cycle positioned after a clock cycle for judging a condition, is only a branch sub-

instruction and a jump sub-instruction. However, it is possible to add a delay function to arbitrary operation sub-instructions such as a load/store sub-instruction, an arithmetic operation, and a shift arithmetic

5 operation. The six registers which are further added to each operation unit have completely the same function as that of the register 33 of the PC controlling part 13 and hold a description indicating a time for executing an operation to be executed in each operation unit.

10 Further, the six memory circuits which are further added to each operation unit have completely the same function as these of the memory circuit 34 of the PC controlling part 13 and hold a description of a result of judging an execution condition of the operation to be executed in
15 each operation unit.

For example, so constituted microprocessor can prepare a sub-instruction as follows:

DADD F0T #H20, H30 R5, R5, R6

20

This sub-instruction is an add sub-instruction which adds a content of the register R5 to a content of the register R6 under a condition that the flag F0 is true and writes a result of the addition to the register R5.

25 This add sub-instruction has two offset values '20' and '30' as immediate values. If an address of the add sub-instruction DADD is No. 2000, the condition of the add

sub-instruction DADD is judged at a clock cycle at which the instruction execution stage E/M of an operation sub-instruction positioned at address No. $2000 + 20 = 2020$ is processed. And, the add operation is executed at a
5 clock cycle at which the execution stage E/M of an operation instruction of the address No. $2000 + 30 = 2030$ is processed. If the add sub-instruction DADD is operation_0 in Figure 2, the addition is conducted by judging the condition in ALU 15; and if it is
10 operation_1, the addition is conducted by judging the condition in ALU 19.

Needless to say, in this embodiment 2, although a register and a memory circuit are added to each operation unit, these can be provided in only selected
15 operation units among the plurality of the operation units. In such a case, a delay function is added to only operation sub-instructions by which the operation units having these registers and memory circuits can process.

As described in the above, the first advantage of a
20 data processing device according to the present invention is that a degree of freedom in scheduling instructions is increased by enabling a second operation instruction to be described prior to another operation instruction for determining an execution condition of
25 the second operation instruction in a program sequence, because a first operation instruction is decoded in a first period and executed in a second period succeeding

thereto, while the second operation instruction of which operation is executed under a predetermined condition is decoded in a third period and executed by judging the condition after passing at least the same time as the
5 second period from the ending of the third period in a fourth period, wherein even though a conditional operation instruction is decoded, the instruction is executed by delaying a time for judging the execution condition not like the first operation instruction which
10 is executed immediately, whereby it is possible to execute the operation instruction which determines the execution condition of the second operation instruction during the delay.

The second advantage of a data processing device
15 according to the present invention is that a degree of freedom in scheduling instructions can further be increased because a time for executing the operation designated by the second operation instruction can further be delayed from the time for judging the
20 condition, whereby the conditional operation instruction can be appropriately executed even though a value which determines the condition of the second operation instruction is overwritten by processings of other operation instructions during the delay.

25 The third advantages of a data processing device according to the present invention is that a degree of freedom in scheduling instructions is increased because

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a first register through a third register which are used
at a time of processing a conditional operation
instruction are provided in an instruction execution
unit, wherein the first register holds a control signal
5 outputted from an instruction decoder; the second
register holds a first description indicating an
execution condition of an operation designated by the
conditional operation instruction; the third register
holds a second description indicating a time for
10 starting judgement of the condition; and the operation
instruction is executed by judging the condition after a
predetermined time is elapsed from the decoding of the
operation instruction based on these three descriptions,
whereby an operation instruction which determines an
15 execution condition of a second operation instruction is
executed during the delay in scheduling the
instructions.

Obviously, numerous modifications and variations of
the present invention are possible in light of the above
20 teachings. It is therefore to be understood that within
the scope of the appended claims, the invention may be
practiced otherwise than as specifically described
herein.